

Multicoherences

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Workshop in Honour of Thomas Ehrhard's 60th Birthday

Outline

Power coherence spaces

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Uniformity

Neutral web

Accuracy

Coherent Differential PCF

Twenty years old story





Power coherence spaces

Definition

A **power** is an endofunctor of the category of sets and inclusions.

Definition

Let P be a power, a **P-coherence space** X is a triple $(|X|, \circ_X, \succ_X)$, $|X|$ the web at most countable, $\circ_X, \succ_X \subseteq P(|X|)$ and $\circ_X \cup \succ_X = P(|X|)$.

Examples

- ▶ \emptyset -coherence space : relational semantics, **Rel**
- ▶ $\mathcal{P}_{\text{fin}}^*$ (finite non-empty sets) : hypercoherences, **NHc**
- ▶ $\mathcal{M}_{\{2\}}$ power = pairs : coherence spaces, **NCoh**_{2}
- ▶ $\mathcal{M}_{\mathbb{N} \setminus \{0,1\}}$ = finite multisets $\# \geq 2$, multicoherences **NMc**



Coherence

Definition

Let X be a P -coherence space. In $P(|X|)$, \circ_X is the **coherence relation**, \succ_X is the **incoherence relation**, $P(|X|) \setminus \circ_X$ the **strict incoherence relation**, $P(|X|) \setminus \succ_X$ the **strict coherence relation** and $\circ_X \cap \succ_X$ is the **neutral relation**, denoted N_X .

Definition

A **clique** of X (a P -coherence space) is a $x \subseteq |X|$ s.t. $P(x) \subseteq \circ_X$.
The linear negation exchanges \circ_X and \succ_X .

Definition

X is **deterministic** if for each clique x of X , each clique y of X^\perp , $\#(x \cap y) \leq 1$.

Uniformity

Definition

(Informal) an expectation to interact with a fair value.

Example

$$M = \lambda b. \quad \text{if } b \quad \text{then } \{\text{if } b \text{ then } M_1 \text{ else } M_2\} \\ \quad \quad \quad \text{else } \{\text{if } b \text{ then } M_3 \text{ else } M_4\}$$

- Multiset-based exponentials:

$$\{([t, t], v_1), ([t, f], v_2), ([t, f], v_3), ([f, f], v_4)\}$$

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- Multiset-based exponentials:

$$\{([t, t], v_1), ([t, f], v_2), ([t, f], v_3), ([f, f], v_4)\}$$

- Set-based exponentials: $\{(\{t\}, v_1), (\{f\}, v_4)\}$

Neutral web

Property

NCoh, **NHc**, **NMc** enjoy a stronger¹ property than determinism :
 $N_X \subseteq \cup_{a \in |X|} P(\{a\})$.

Definitions

The **neutral web** $|X|_N$ is the set $\{a \in |X| \mid P(\{a\}) \subseteq N_X\}$. The **neutral restriction** of X is the sub-space of X of web $|X|_N$, that is $(|X|_N, \circ_X \cap M, \asymp_X \cap M)$ where $M = P(|X|_N)$, and the neutral restriction of a clique x of X is $x \cap |X|_N$.

The space is reduced to the expected sites of interactions.

¹provided that the power is strictly monotone and preserves disjointness

Extensional collapse

Definition

The **extensional collapse** of a semantics, with respect to the simple types hierarchy, is its quotient by the partial equivalence relations \sim_A , given by equality on basis types and :

$$f \sim_{A \Rightarrow B} g \text{ iff if } x \sim_A y \text{ then } f(x) \sim_B g(y).$$

Theorem

*The extensional collapse of **NCoh** and of the multiset-based **Coh** semantics is the set-based **Coh** semantics. The same for **NHc/Hc**, **NMc/Mc**.*

The neutral restriction plays a crucial role in the proof (all the commutations one can expect).

Accuracy

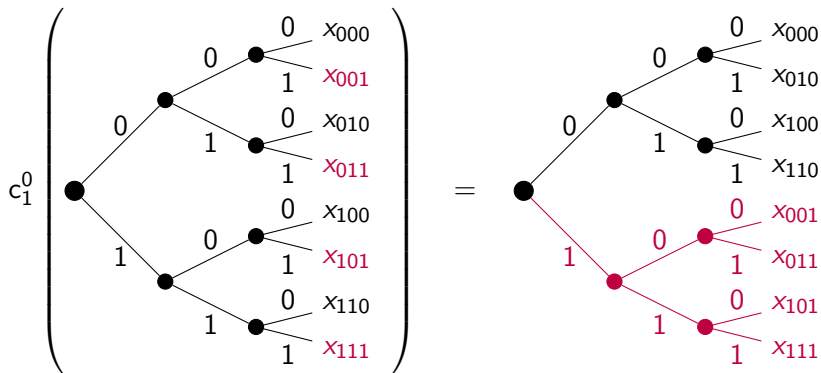
- ▶ No full completeness (too simple?)
- ▶ Gustave function in coherence spaces
- ▶ In **Hc**, **Mc** subdefinability: for any clique x of $\mathbf{Bool}^n \rightarrow \mathbf{Bool}$ there exists a term M s.t. $x \subseteq [M]$.
- ▶ **Hc** extensional collapse of the sequential algorithms
- ▶ non-uniform is non-sequential. $\{([t], t), ([f], t), ([t, f], t)\}$

Coherent Differential PCF syntax

- ▶ Thomas Ehrhard, A coherent differential PCF. 2022
- ▶ Differential lambda-calculus without the general sum $M + N$
- ▶ **NCoh**, **NHc**, **NMc** are valid semantics
- ▶ types: $A, B := D^n \iota \mid A \Rightarrow B$
- ▶ D applies to any type but $D(A \Rightarrow B) = A \Rightarrow DB$
- ▶ D applies to terms and triggers a rewriting
 $D(\lambda x^A.M : A \Rightarrow B) = \lambda.x^{DA} \partial(x, M) : DA \Rightarrow DB$
- ▶ $|D^n \iota| = \{(\delta, i) \mid i \in \mathbb{N}, \delta \text{ word of } n \text{ bits}\}$
- ▶ projections π_0, π_1 ; injections in_0, in_1 ; swap c ; delayed sum θ
- ▶ succ^d , if_A^d , let_A^d , π_0^d , $\theta^d \dots$ at **depth** d
- ▶ Don't confuse $D^0 \iota = \iota$ but $\pi_i^0(M) = \pi_i(M)$.



Swap





Rel intersection typing system (examples)

$$\frac{\pi_0}{\frac{\frac{\frac{\frac{\vdots}{x : \iota \vdash 2 : \iota} \text{typing}}{x : [1] : \iota \vdash \text{if}_\iota(x, 2, 3) : 3 : \iota} \text{if1}}{x : [0, 1] : \iota \vdash \text{if}_\iota(x, \text{if}_\iota(x, 2, 3), \text{if}_\iota(x, 4, 5)) : 3 : \iota} \text{abs}}{\vdash \lambda x^\iota. \text{if}_\iota(x, \text{if}_\iota(x, 2, 3), \text{if}_\iota(x, 4, 5)) : ([0, 1], 3) : \iota \Rightarrow \iota} \text{if0}} \pi_2}$$

Where :

$$\pi_0 = \frac{}{x : [0] : \iota \vdash x : 0 : \iota} \text{var} \quad \pi_1 = \frac{}{x : [1] : \iota \vdash x : 1 : \iota} \text{var}$$

$$\pi_2 = \frac{\vdots}{x : \iota \vdash \text{if}_\iota(x, 4, 5) : \iota} \text{typing}$$



Example 2, $x : D_\iota$

$$\frac{\frac{\frac{\pi_0}{x : [(\delta, 0), (\delta', 1)] : D_\iota \vdash \text{if}_\iota^1(x, \text{if}_\iota^1(x, 2, 3), \text{if}_\iota^1(x, 4, 5)) : (\delta\delta', 3) : D^2_\iota} \text{if0}}{\vdash \lambda x^\iota. \text{if}_{D_\iota}^1(x, \text{if}_\iota^1(x, 2, 3), \text{if}_\iota^1(x, 4, 5)) : ([(\delta, 0), (\delta', 1)], (\delta\delta', 3)) : D_\iota \Rightarrow D^2_\iota} \text{abs}}{\frac{\frac{\frac{\pi_1 \quad \pi_3 \quad \frac{\text{num}}{x : [] : D_\iota \vdash 3 : \iota}}{\text{if1}}}{x : [(\delta', 1)] : D_\iota \vdash \text{if}_\iota^1(x, 2, 3) : (\delta', 3) : D_\iota} \text{if1}}{x : [(\delta, 0), (\delta', 1)] : D_\iota \vdash \text{if}_\iota^1(x, \text{if}_\iota^1(x, 2, 3), \text{if}_\iota^1(x, 4, 5)) : (\delta\delta', 3) : D^2_\iota} \text{if0}}{\vdash \lambda x^\iota. \text{if}_{D_\iota}^1(x, \text{if}_\iota^1(x, 2, 3), \text{if}_\iota^1(x, 4, 5)) : ([(\delta, 0), (\delta', 1)], (\delta\delta', 3)) : D_\iota \Rightarrow D^2_\iota} \text{abs}} \text{abs}}$$

Where :

$$\pi_0 = \frac{\text{var}}{x : [(\delta, 0)] : D_\iota \vdash x : (\delta, 0) : D_\iota}$$

$$\pi_1 = \frac{\text{var}}{x : [(\delta', 1)] : D_\iota \vdash x : (\delta', 1) : D_\iota}$$

$$\pi_2 = \frac{\vdots}{x : D_\iota \vdash \text{if}_\iota^1(x, 4, 5) : D_\iota} \text{typing} \quad \pi_3 = \frac{\vdots}{x : D_\iota \vdash 2 : \iota} \text{typing}$$



Example 3, differentiation

$$\frac{\frac{\frac{\pi_0}{x : [t, f] : \iota \vdash \text{if}_\iota(x, \text{if}_\iota(x, 2, 3), \text{if}_\iota(x, 4, 5)) : 3 : \iota} \quad \frac{\frac{\pi_1}{x : [f] : \iota \vdash \text{if}_\iota(x, 2, 3) : 3 : \iota} \quad \frac{\pi_3}{x : [] : \iota \vdash 1 : \iota} \text{ num}}{\text{if1}}}{\text{if0}} \quad \pi_2}{\vdash \lambda x^\iota. \text{if}_\iota(x, \text{if}_\iota(x, 2, 3), \text{if}_\iota(x, 4, 5)) : ([t, f], 3) : \iota \Rightarrow \iota} \text{ abs}}{\vdash D(\lambda x^\iota. \text{if}_\iota(x, \text{if}_\iota(x, 2, 3), \text{if}_\iota(x, 4, 5))) : (([0, t], (1, f)], (1, 3)) : \iota \Rightarrow \iota} \text{ map} \quad a \in \partial_{[[\iota]]}$$

Where $a = (([0, t], (1, f)], (1, [t, f]))$ and :

$$\pi_0 = \frac{}{x : [t] : \iota \vdash x : t : \iota} \text{ var} \quad \pi_1 = \frac{}{x : [f] : \iota \vdash x : f : \iota} \text{ var}$$

$$\pi_2 = \frac{\vdots}{x : \iota \vdash \text{if}_\iota(x, 4, 5) : \iota} \text{ typing} \quad \pi_3 = \frac{\vdots}{x : \iota \vdash 2 : \iota} \text{ typing}$$

Conclusion

- ▶ Uniformity appeared to be a kind of anticipation of the sites of interaction
- ▶ This intuition works pretty well for the closed case where a clique encounter an anti-clique, as for the extensional collapses.
- ▶ I am still digging to understand coherent differential PCF.
- ▶ This calculus admits **NCoh**, **NHc**, **NMc** but ~~not their uniform versions~~².
- ▶ What is wrong? Is there an answer in terms of interaction?

²That was a mistake in my talk, as confirmed by T. Ehrhard, uniform coherent spaces are a model of Coh. Diff. PCF

Conclusion

Thank you!

